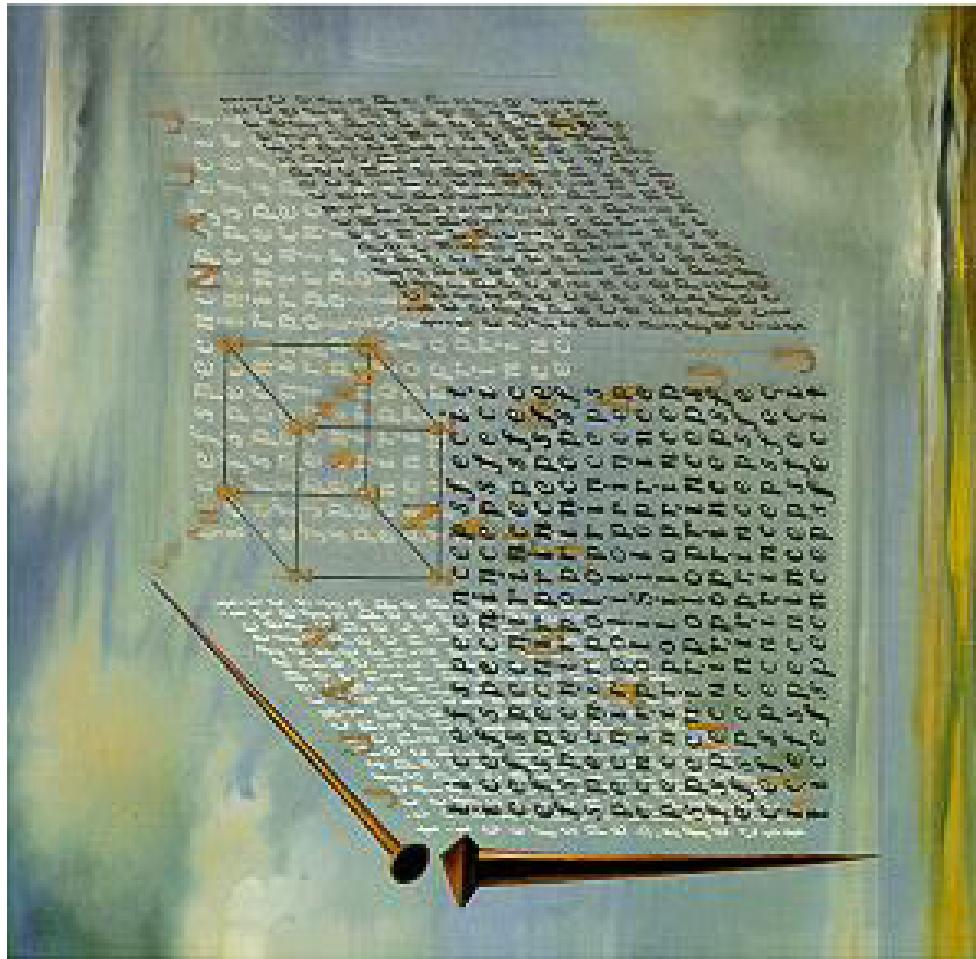


# Raster Databases

- tutorial -

VLDB 2007  
Vienna, 25-sep-2007

Peter Baumann  
**Jacobs University Bremen,  
rasdaman GmbH**



# About the Presenter

[www.faculty.jacobs-university.de/pbaumann](http://www.faculty.jacobs-university.de/pbaumann)

## ■ Professor of Computer Science

- research focus: large-scale multi-dimensional raster services
  - ...and application in geo, life science, Grid, and e-learning
- geo raster service standardization: OGC
- research spin-off: rasdaman GmbH

## ■ Jacobs University Bremen

- Private research university, est. 1998 by State of Bremen
- >1100 Studenten, 91 nations, 25% German
  - ACQUIN accredited
  - Transdisciplinary, international, multi-cultural, all-english
- "Smart Systems" CS graduate program
  - MSc, PhD



# Roadmap

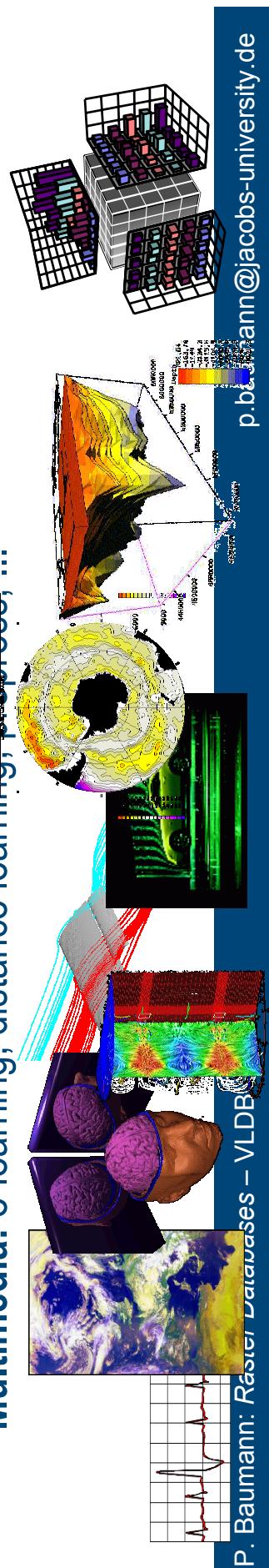


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- **Introduction**
- **Conceptual modelling**
- **Architecture**
  - Arch I: Storage Management
  - Arch II: Query Processing
- **Applications**
- **Wrap-up**

# Why (Large) Arrays?

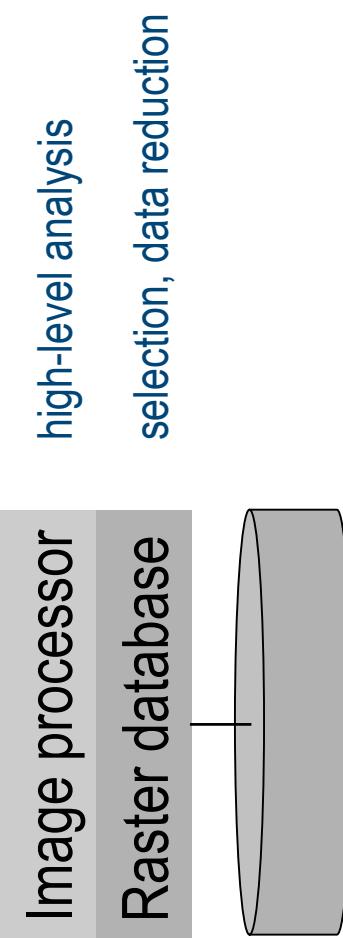
- Key characteristics: **Dimensional, gridded (Euclidean space), large**
  - raster = array = Multidimensional Discrete Data (MDD)
- **Sensor, image, statistics data**
  - **Life Science:** Pharma/chem, healthcare / bio research, bio statistics, genetics
  - **Geo:** Geodesy, geology, hydro/ocean, meteorology, earth system research, ...
  - **Management/Controlling:** statistics / Decision Support, OLAP, Warehousing, ...
  - **Engineering & research:** **Simulation** & experimental data in automotive/shipbuilding/aerospace industry, turbines, process industry, **astronomy**, experimental physics, high energy physics, ...
  - **Multimedia:** e-learning, distance learning, prepress, ...



# Raster Services: Differentiation



- multimedia databases
  - Analyse images, then drop them and **work on auxiliary structure**
- image processing
  - Advanced processing of rasters, but not on objects **>> main memory size**
- image understanding,  
computer vision
  - General recognition **probabilistic**
  - databases to deliver **exact results** whenever possible
- Statistical DB / OLAP: dense vs sparse

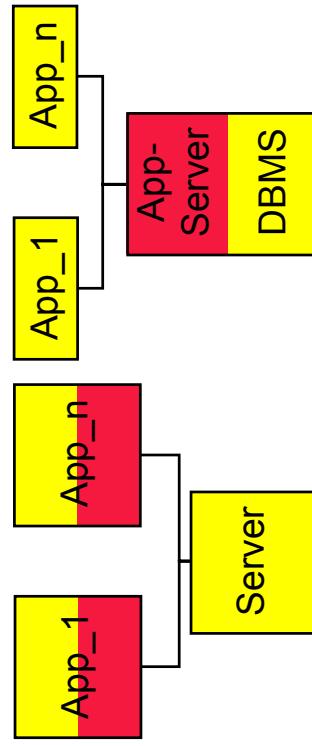


# Why Array Databases?



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- Why should we bother?  
*...because it's tons of data, that's us!*
  - Multi-Terabyte objects, soon multi-Petabyte archives
- What can we offer?  
*...,"Classical" database benefits, for a new data type:*
  - information integration
  - flexibility
  - scalability
  - ..plus all our further assets



# Roadmap



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- Architecture
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  - Arch II: Query Processing
- Applications
- Wrap-up

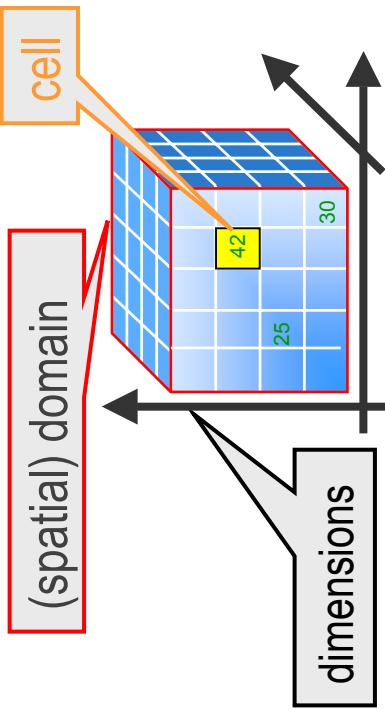
# History

- Database view on raster images (eg, [XXX]):
  - „image data...*matrix of pixels*“, but: „data appear just as a *string of bits*“ → BLOBs
- Steps towards array support:
  - Image partitioning (tiling) in standardised files, API access library [Tamura 1980]
  - Fixed set of imaging operators (scaling, rotation, edge extraction, thresholding, ...)  
[Chang, Fu 1980; Stucky, Menzi 1989; Neumann et al 1992]
  - PICDMs [Chock, Cardenas 1984]: image stack (same res); no nesting; no architecture  
rasdaman array algebra [Baumann 1991] & system [Baumann 1994+]
  - AQL [Libkin, Machlin, Wong 1996; Machlin 2007]
  - AML [Marathe & Salem 1997, 1999]; RAM [Ballegooyij, de Vries, Kersten 2003];  
[Ordinez, Garcia 2007]
  - ESRI ArcSDE, Oracle GeoRaster [200x]

# Conceptual Modelling: Array Algebra

- **Array = function:**

- $a: X \rightarrow F$ ,  $a = \{ (x, f) : x \in X, a(x) = f \in F \}$   
for finite multi-dimensional interval  $X \subset \mathbb{Z}^d$ ,  $d > 0$ , algebraic structure  $F$
- $d$ : Dimensionality of  $a$ ,  
 $X$ : spatial domain,  
 $F$ : Value set (range), Pixel, Voxel, ...



- 3 primitives:

- Array constructor
- Condenser
- Sort

- Inspired by AFATL Image Algebra [Ritter et al 1990], basis for rasdaman system

# Array Operations: MARRAY



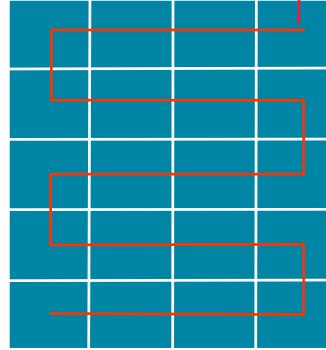
- **Array constructor:**  $\text{MARRAY}_{X,p}( e(p) ) := \{ (p,f) : f = e(p), p \in X \}$ 
  - for n-D finite interval  $X$ , expression  $e(p)$  potentially containing occurrences of  $p$ , of result type  $F$ 
    - Ex:  $\text{MARRAY}_{X,p}( a[p] + b[p] ) =: a + b$   
 $\text{MARRAY}_{X,p}( p[0] )$
- **Shorthand: "induced operations"**
  - $(X = \text{sdom}(a) = \text{sdom}(b), a:X \rightarrow F, b:X \rightarrow G \text{ and } f:F \rightarrow F', g:F \times G \rightarrow G')$ :
    - $f_{\text{ind}} : X^F \rightarrow X^{F'}$ ,  $f_{\text{ind}}(a) = \text{MARRAY}_{X,X}( f(a(x)) )$  unary induced operation
    - $g_{\text{ind}} : X^F \times X^G \rightarrow X^{G'}$ ,  $g_{\text{ind}}(a,b) = \text{MARRAY}_{X,X}( g(a(x), b(x)) )$  binary induced operation

# Array Operations: COND



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- Condenser:  $\text{COND}_{o,X,x}( e(a,x) ) := e(a,p_1) \circ e(a,p_2) \circ \dots \circ e(a,p_n)$ 
  - n-D finite interval  $X$ ,  $\circ$  commutative, associative, **associative**,  $e(a, p)$  expression potentially containing  $a$  and  $p_i$
  - Ex:  $\text{add\_cells}(a) := \text{COND}_{+, \text{sdom}(a), p}( a[p] )$
- Shorthands:
  - `count_cells()`, `avg_cells()`,  
`max_cells()`, `min_cells()`,  
`some_cells()`, `all_cells()`
  - cf. Relational aggregates



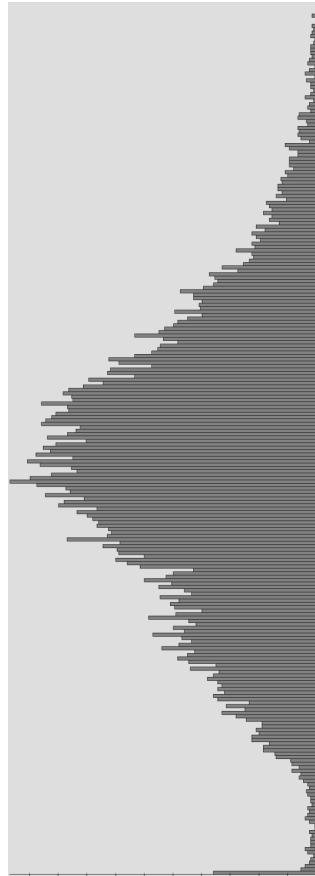
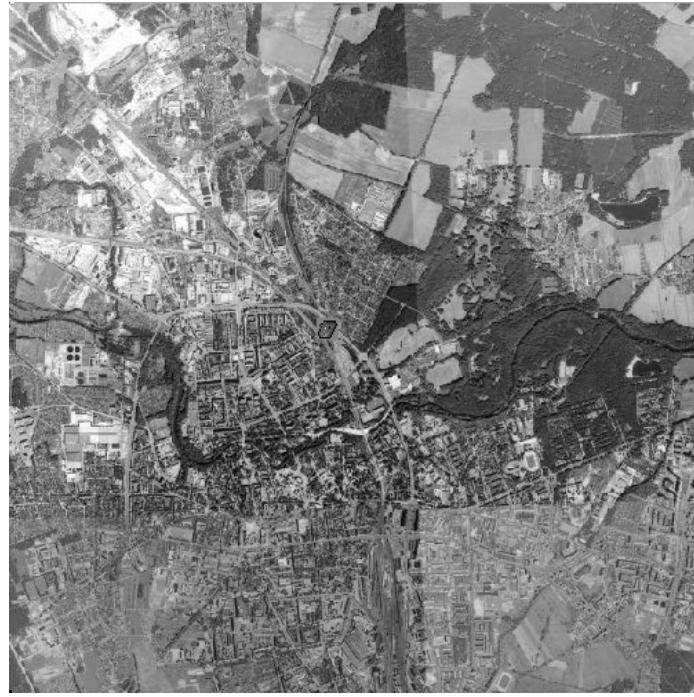
# Example: Histogram

- Histogram of an n-D array over 8-bit unsigned integer:

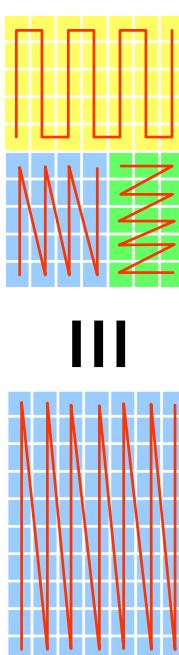
- $\bullet \quad H(a) = MARRAY_a, [0:255] (\text{ count\_cell } s (a = n))$

- MARRAY can change cell type, dimension, domain!

- $\bullet \quad \text{sdom}(H(\text{image})) = [0:255]$



# Properties

- **Array Algebra declarative** wrt array addressing
  - MARRAY: implicit iteration; COND: associative + commutative aggregator functions
  - tile-based processing:  

- **Array algebra safe** in evaluation
  - Array indexing without recursion
  - [Machlin 2007] goes beyond
- **Expressive power:** AML, Array Algebra **equal to relational + ranking** [Libkin, Machlin, Wong 1996]
  - In practice: filters, convolutions, statistics, ...

# From Algebra To Query Language



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- rasdaman ("raster data manager") middleware
  - in commercial use since 2001 (e.g. IGN-F: 13 TB ortho image, PostgreSQL)
- Data model: collections of typed arrays + OIDs
- **Data definition language:** rasdl [ODMG ODL]
  - Parametrised array constructor
  - Ex:

```
typedef marray
< unsigned char, [ 1:1024 , 1:768 ]
> XgaGreyImage;
```
- **Retrieval & manipulation language:** rasql, based on SQL92
  - Select, insert, update, delete; speciality: partial update
  - Set oriented: all queries return sets, ...ahem: multi-sets, ...ahem: lists of arrays

# Inset: Types vs Type Constructors

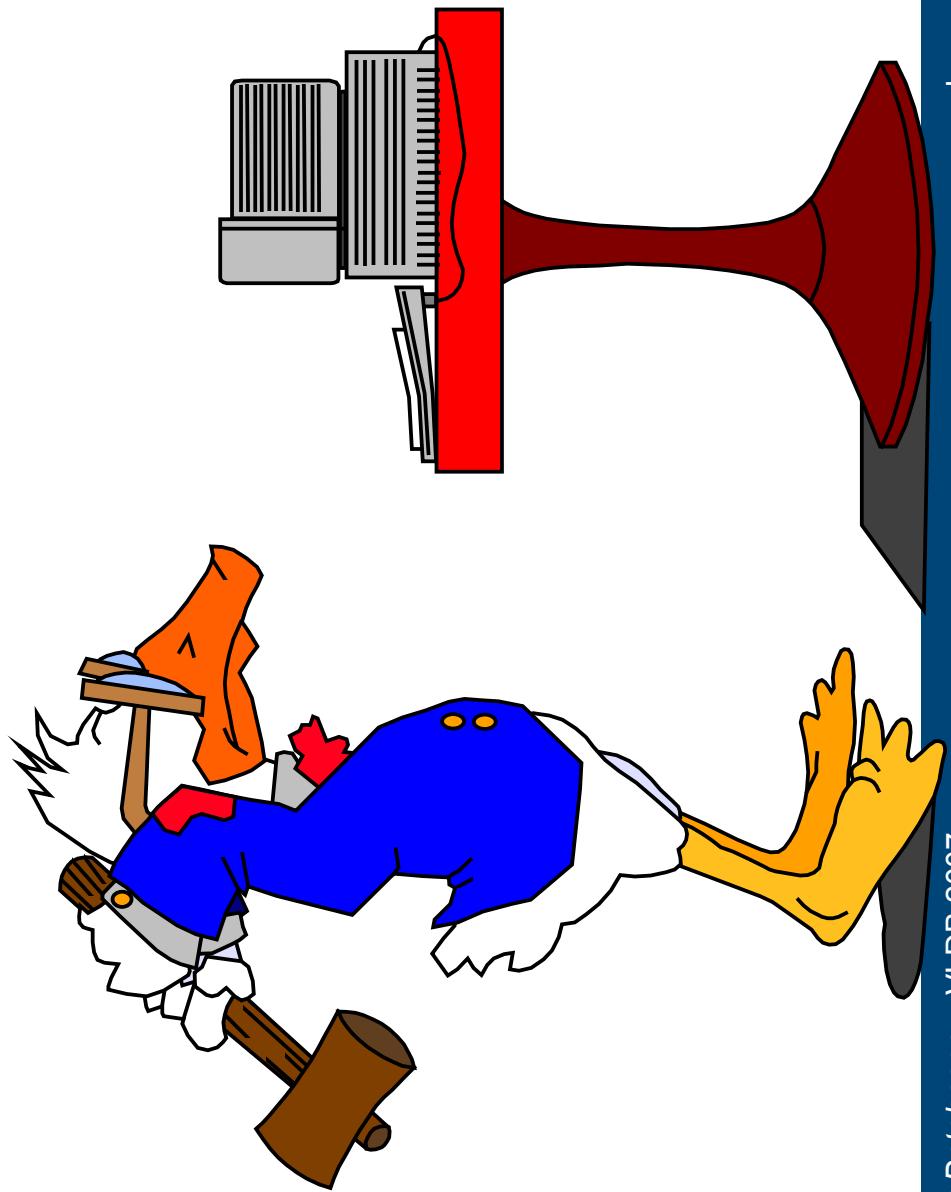
- Remember: Marray is not a *type*, but a parametrized **type constructor**
  - Ex:  

```
typedef marray
< struct { double vx, vy; }, [ 0:*, 0:127, 0:63, 0:16 ]>
ECHAM_T42_Windspeed;
```
  - Cf. Stack: stack<int> is constructor, stack<int> a concrete type
- Object-relational extensions allow user-defined data types,  
however **not** type constructors
  - Exception: Predator, U of Wisconsin-Madison

# Demo



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# Oracle 10g/11g GeoRaster



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- **GeoRaster**
  - Large 2-D geo raster images
  - Response to ESRI's ArcSDE 8
- **Functionality:**
  - (non-transparent) image pyramids
  - Subsetting, component extraction
  - reprojection?
- **Observations**
  - data independence?  
eg, pyramids visible
  - No SQL-integrated processing
    - No optimization found

```
declare
  g sdo_georaster;
  b blob;
begin
  select raster into g
  from uk_rasters
  where id = 4;
  dbms_lob.createTemporary(b, true);
  sdo_georaster.getRasterSubset(
    georaster => g,
    pyramidLevel => 0,
    window =>
      sdo_number_array(0, 0, 699, 899),
      bandnumbers => '0',
      rasterblob => b );
end;
```

```
select g.green[0:699, 0:899]
  from uk_rasters as g
 where oid(g) = 4
```

# Roadmap



- Introduction
- Conceptual modelling
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  - Arch I: Storage Management
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# Storage Mapping



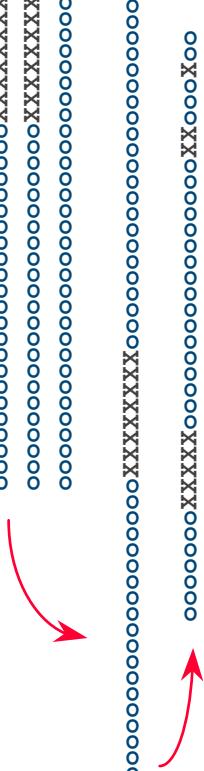
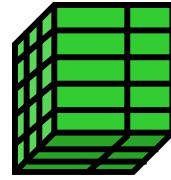
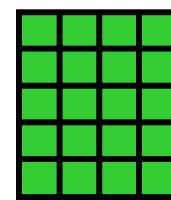
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- Task: materialise finite interval  $X \subset \mathbf{Z}^n$ , find suitable (disk) access structure
  - Core structural property: **Euclidean neighbourhood** in  $\mathbf{Z}^n$
  - Secondary, contents/app based: data density/ sparsity, data pattern, access pattern
- Excursion: difference to arrays in main memory
  - Ex: APL [Iverson 1968]
  - **Assumption 1:**  
access times independent from array position
    - $\text{cost}(\text{,} \alpha [x]) = \text{const for all } ,x$
  - **Assumption 2:**  
access times independent from access sequence
    - $\text{cost}(\text{,} \alpha [x] ; \alpha [y]) = 2 * \text{cost}(\text{,} \alpha [x]) = \text{const for all } ,x, ,y$

# Storage Mapping: Variants

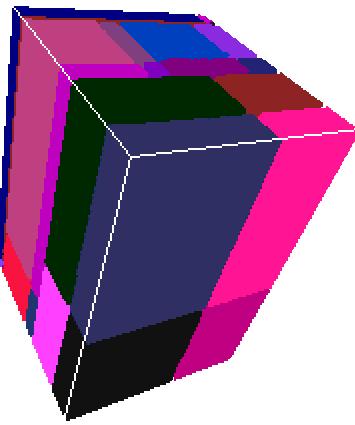
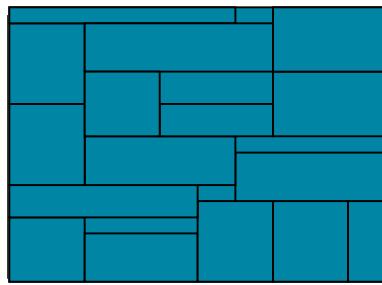


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- BLOB (binary large object)
    - Coordinate free sequence
    - Costs mainly position/dimension dependent
  - 
  - Sequence independent, coordinates explicit  $\{ (x_1, f_1), (x_2, f_2), \dots, (x_n, f_n) \}$ 
    - Costs not position correlated, but high
    - Sequence independent, coordinates explicit
  - Imaging, multidimensional OLAP
    - Partitioning, sequence within partition
    - Costs low for bulk access, usually not location correlated
- 
- 

# Partitioned Array Storage

[Furtado 2000, Widmann 2001]



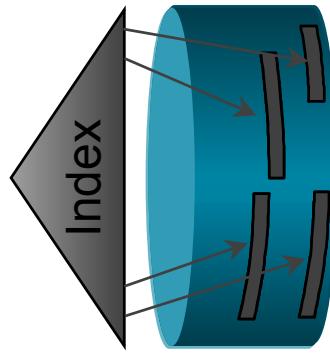
- multidimensional object  
    ⇒ multidimensional tiles

- Tile = subarray
- Also called "chunking"  
[Sarawagi, DeWitt]



- Tiles stored as BLOB  
in relational database

- Compression
- Geo index



# Storage Layout: Tuning



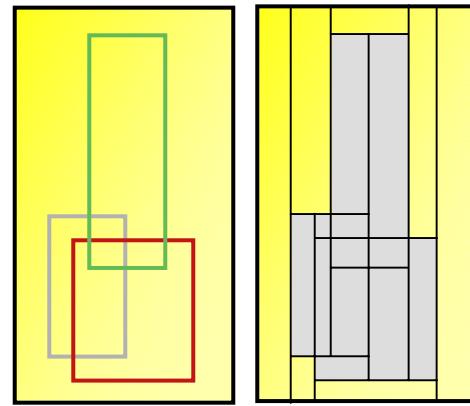
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- Parameters:
  - Tiling strategy
  - Geo index
  - Data format within tiles, incl compression
- Many dependencies
  - Access patterns, data contents
  - Buffer size, page size, CPU performance, bus bandwidth, ...
- In rasdaman:
  - Controlled via API,  
eg `rasj class RasSTorageLayout`
  - Storage layout determined during insertion
  - Reorganisation = copying (beware!),  
possible via API

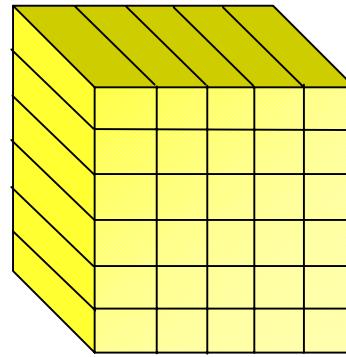
# Tiling Strategies

- Goal: **faster tile loading** by adapting storage units to access pattern
- Issues
  - When is tiling optimal? Tiling strategies?
- 3 sample tiling strategies [Furtado 1999]:

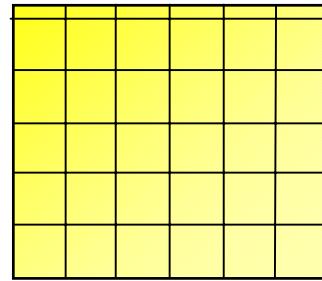
area of interest



directional



regular



# Tile Based Compression

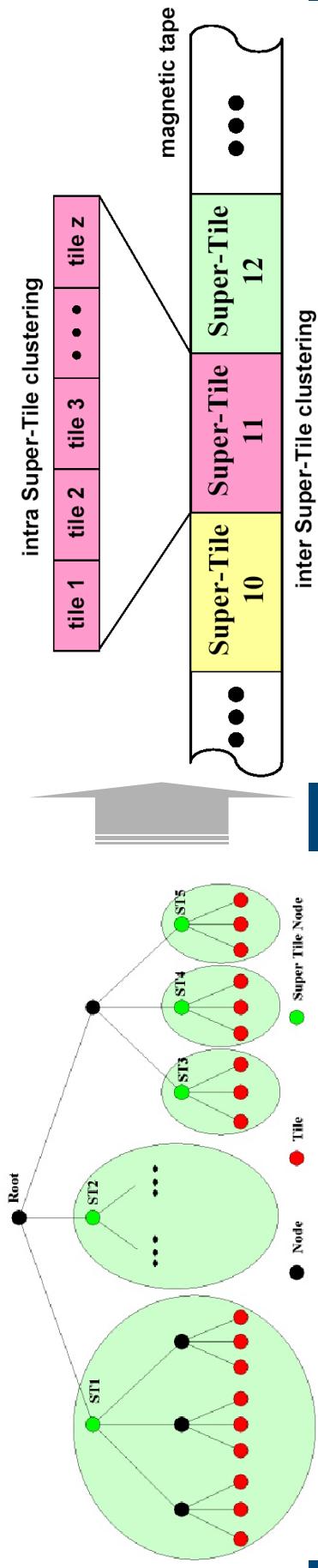


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- Starting point: tiles are unit of access → compression units are tiles, too
- Degree of freedom (tuning parameter):  
mixing of compression methods according to **access pattern**
  - uncompressed ← hot spots
  - Fast & less storage gain ← high volume, frequent access
  - Slow & high storage gain ← infrequent access, high volume

# Hierarchical Storage Media Management

- near-line tape archives as storage extensions  
[Sarawagi, Stonebraker 1994]
- Issue: respect spatial clustering
  - access locality, long positioning times!
- **super tile** = tile set under some index node  
[Reiner 2001]
  - Natural unit, comfortable to handle (eviction information in index node!)



# Roadmap

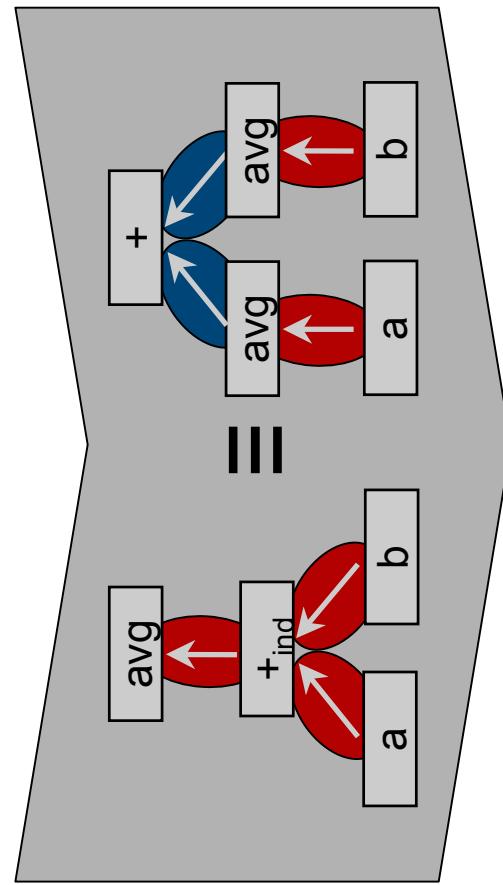


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- Introduction
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# Query Optimization

```
select avg_cells( a + b )
from a, b
```



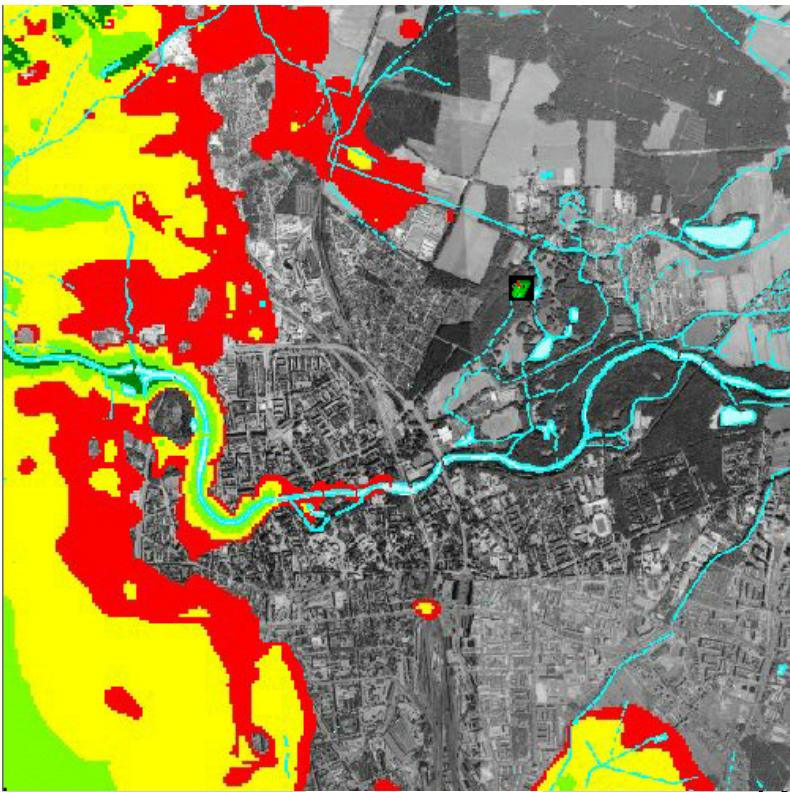
```
select avg_cells( a )
+ avg_cells( b )
from a, b
```

- *understood:*
  - heuristic optimization
  - partially understood:
    - cost-based optimization

# Optimisation Does Pay Off!

- Complex queries give more space to optimizer
- Typical OGC Web Map Service query:

```
select jpeg(  
    scale(bild0[...],[1:300,1:300]) * {  
        overlay ((scale(bild1[...],[1:300,1:300])<71.0)) *  
        overlay bit(scale(bild2[...],[1:300,1:300]), 2) * {  
            overlay bit(scale(bild2[...],[1:300,1:300]), 5) * {  
                overlay bit(scale(bild2[...],[1:300,1:300]), 7) * {  
                    overlay bit(scale(bild2[...],[1:300,1:300]), 6) * {  
                        overlay bit(scale(bild2[...],[1:300,1:300]), 3) * {  
                            overlay bit(scale(bild2[...],[1:300,1:300]), 4) * {  
                                overlay bit(scale(bild2[...],[1:300,1:300]), 1) * {  
                                    overlay bit(scale(bild2[...],[1:300,1:300]), 0) * {  
                                        {102c, 102c, 102c}  
                                        {255c, 255c, 0c}  
                                        {191c, 242c, 128c}  
                                        {191c, 255c, 255c}  
                                        {0c, 255c, 255c}  
                                        {102c, 102c, 102c}  
                                    )  
                                )  
                            )  
                        )  
                    )  
                )  
            )  
        )  
    )  
}
```

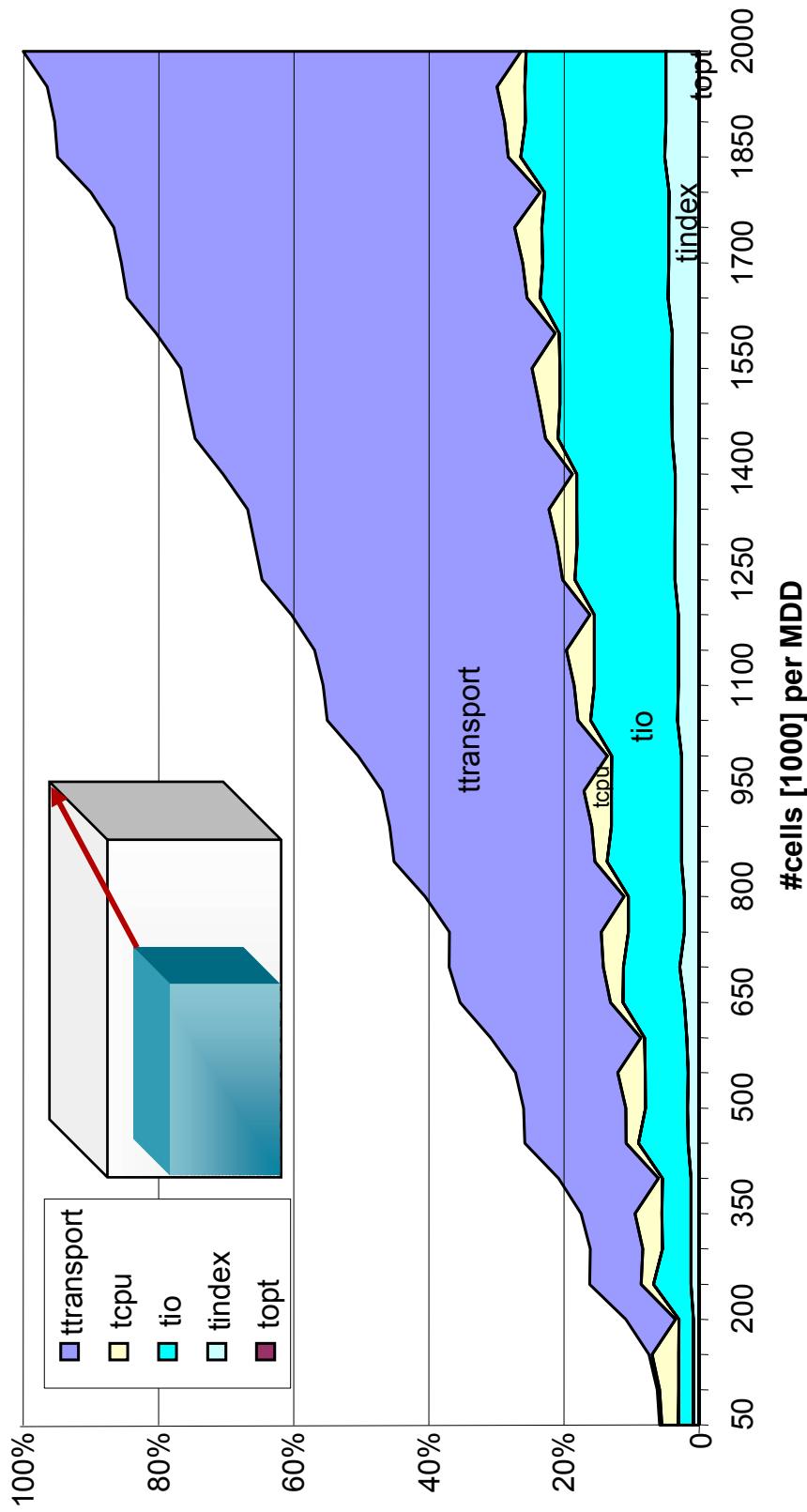


# Benchmarks: Data Access

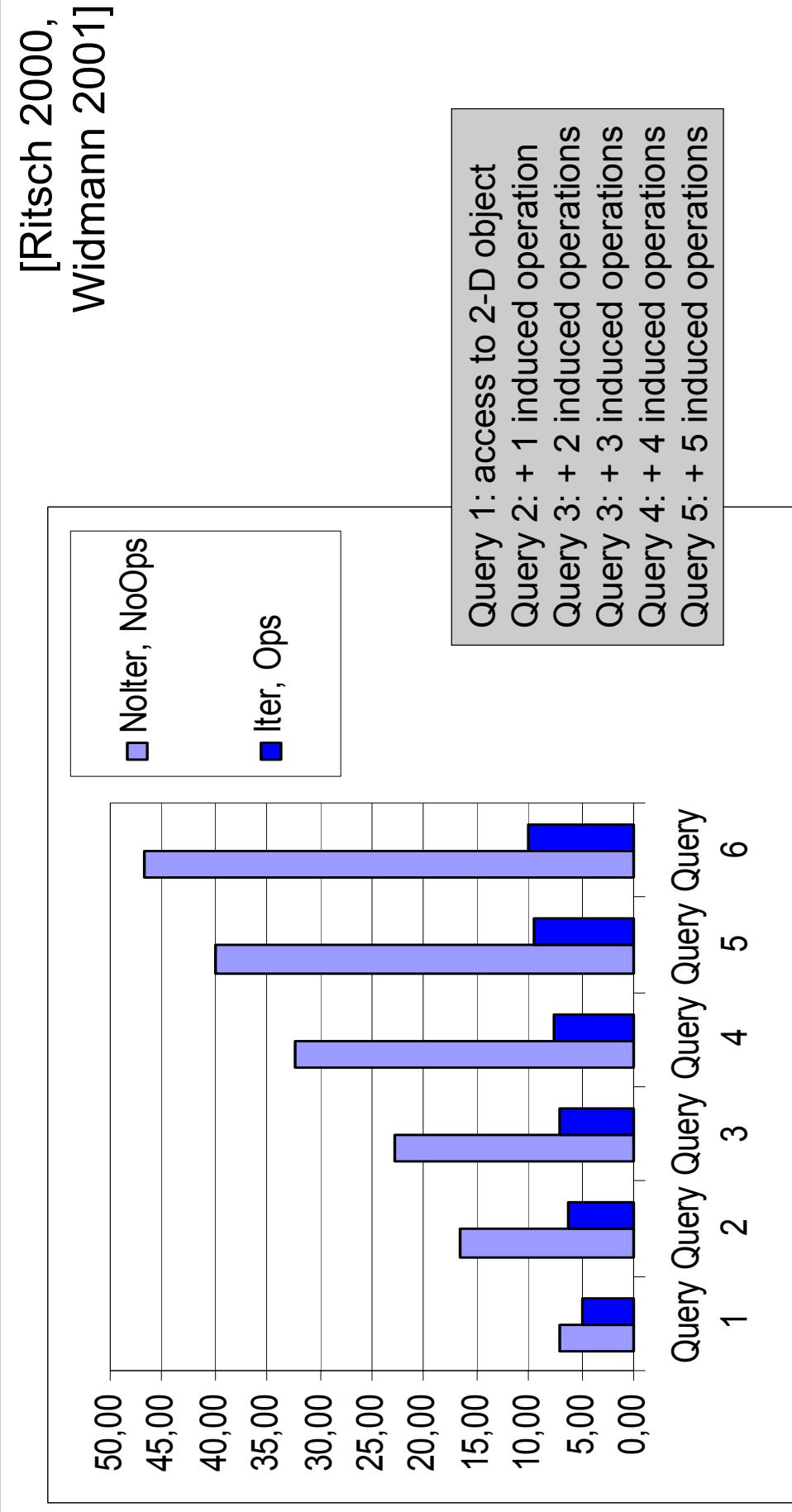


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[Ritsch 2000, Widmann 2001]

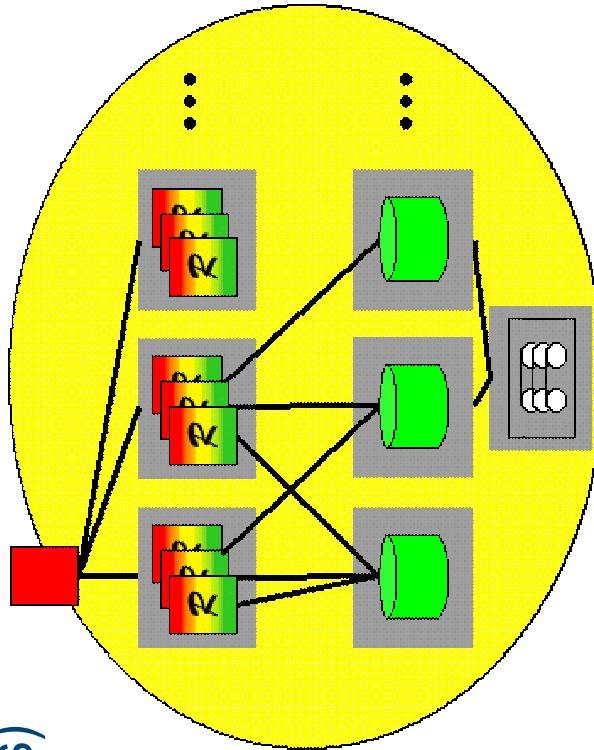


# Benchmarks: Data Processing



# Query Parallelisation

- easy: inter-query parallelization  
(one client – one dedicated server process)
  - Long-runners don't block service
    - higher throughput
- Non-trivial: intra-query parallelization  
(one client – several server processes)  
[Hahn 2003]
  - Idea: tiles dynamically assigned to processors
  - *Non-trivial array index patterns?*

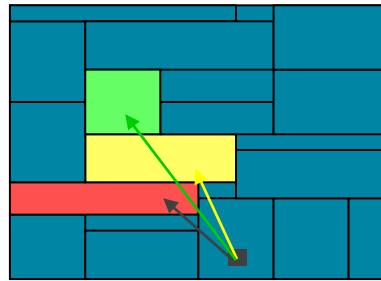


# Non-Local Access Patterns

- **Problem:** how to efficiently evaluate tiles

in face of non-trivial access patterns

- marray  $x$  in  $X$   
values  $\text{img}[\text{ } \underline{f(x)} \text{ }]$



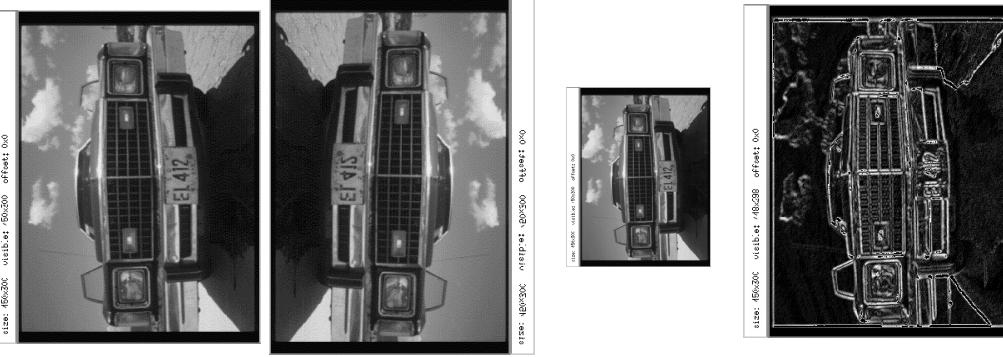
- **Use cases:**

- mirroring:  $f(x) = h_i - x$
- scaling:  $f(x) = x / s$
- Filtering: marray  $x$  in  $\text{sdom}(\text{img})$

```
values condense +  
over y in  $\text{sdom}(\text{kernel})$   
using  $a[x+y] * \text{kernel}[y]$ 
```

1	3	1
0	0	0
-1	-3	-1

- **Approach:** address important cases first: const, linear expressions



# Roadmap



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# Geo Service Standardization



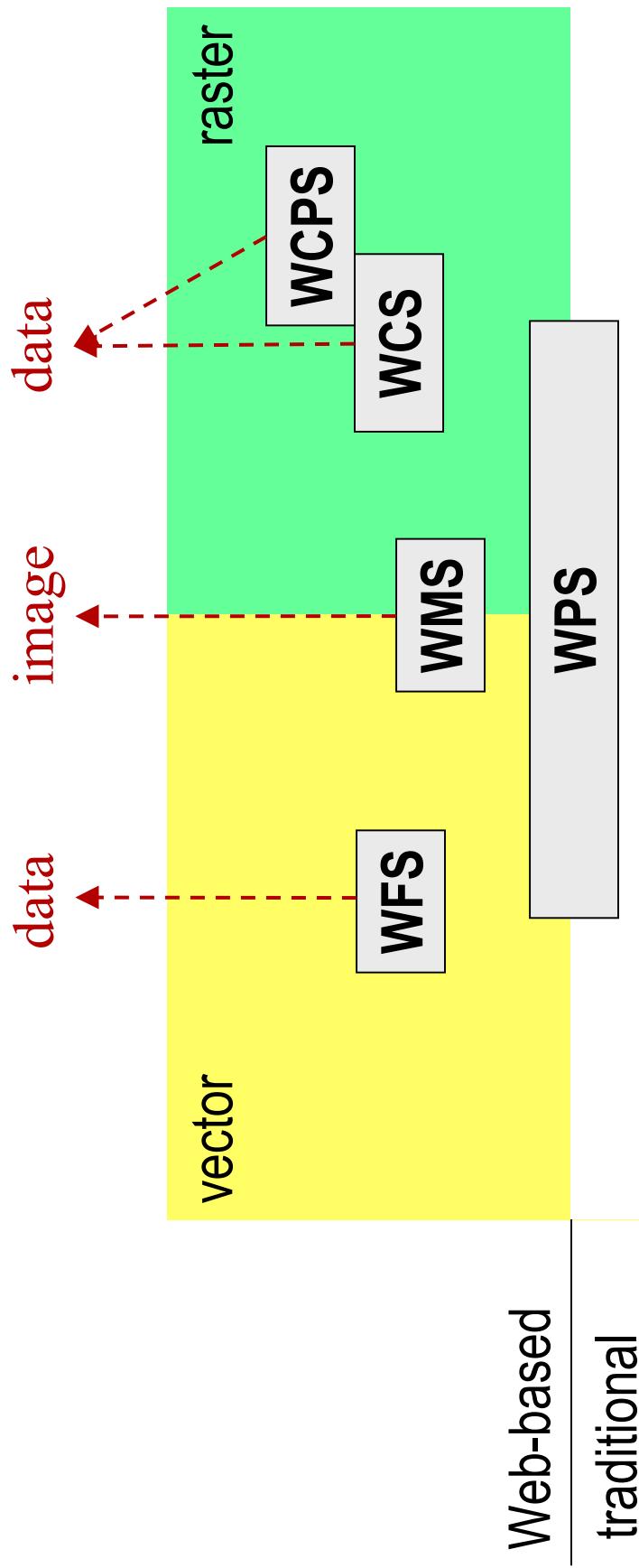
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- OGC (Open GeoSpatial Consortium) driving geo service standards
  - Web-based modular, open, interoperable geo services
  - Liaisons with ISO TC 211, OASIS, CGI/IUGS; ...
  - [www.opengeospatial.org](http://www.opengeospatial.org)
- Raster = coverage in OGC / GIS speak
  - Web Coverage Service Revision Working Group (WCS.RWG)
  - Web Coverage Processing Service Group (WCPS)
    - Coverages WG
  - GALEON OGCI network (Geo-interface to Atmosphere, Land, Earth, Ocean, NetCDF)

# (Part of) The OGC Quilt



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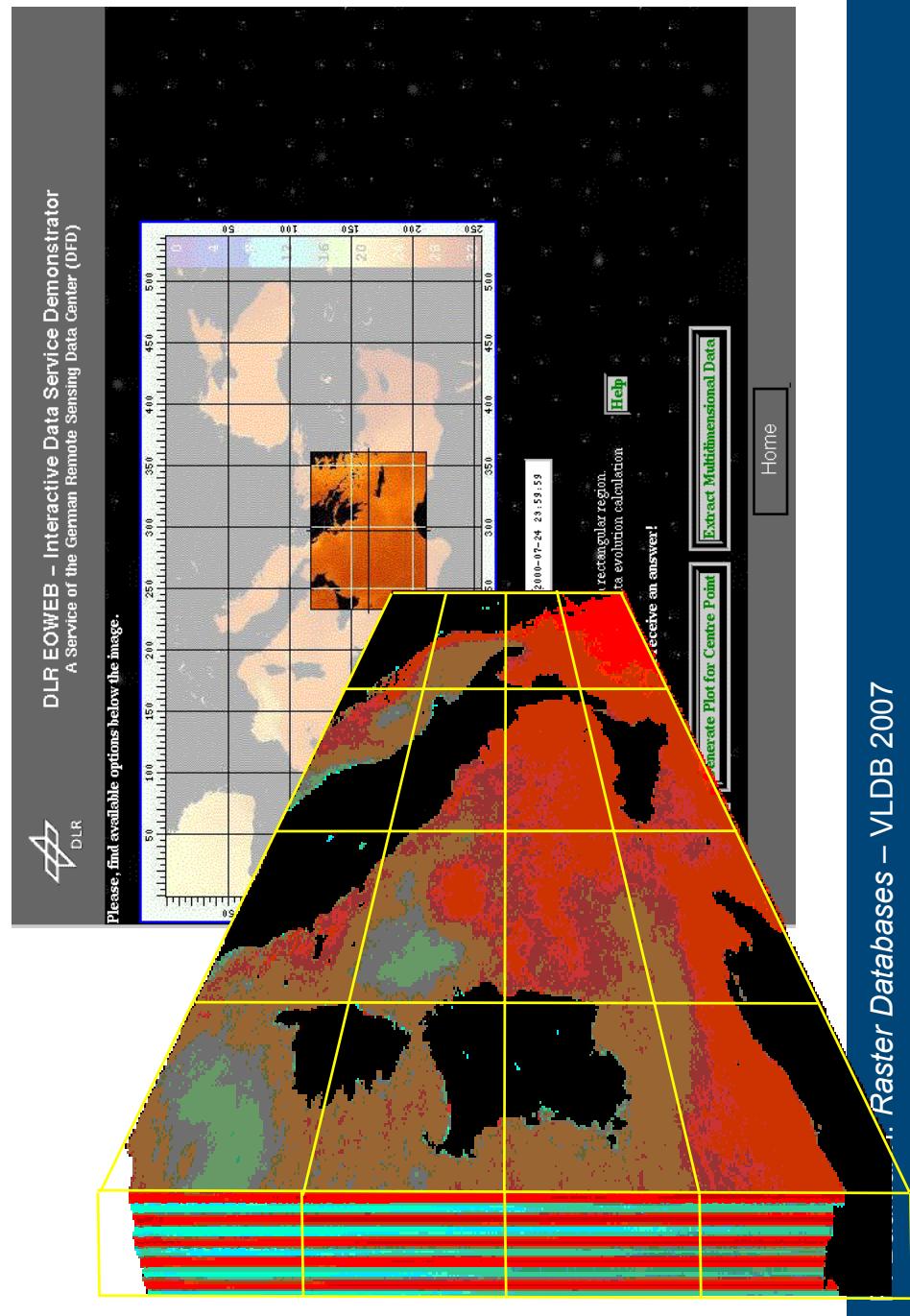
- WMS "portrays spatial data → pictures"
- WCS: "provides data + descriptions; data with *original semantics*, may be interpreted, extrapolated, etc."

# Sample WCS Based 3-D Service



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DLR-DFD: eoweb.dlr.de [Diedrich et al 2001]  
based on rasdaman



# WCPS



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- Request yields one or more n-D coverages

- Abstract syntax (requests shipped as XML):

```
for var in ( coverageList )
[ where condition(var) ]
return processingExpr(var)
```

- Example:

```
for m in ( ModisA, ModisB, ModisC )
where
max( m.red > 127 )
return
encode( m.red + m.nir,
"tiff"
)
```

( tiff\_A,  
tiff\_C )



# Climate Modelling

## ■ Example: ECHAM T42 (cf. video)

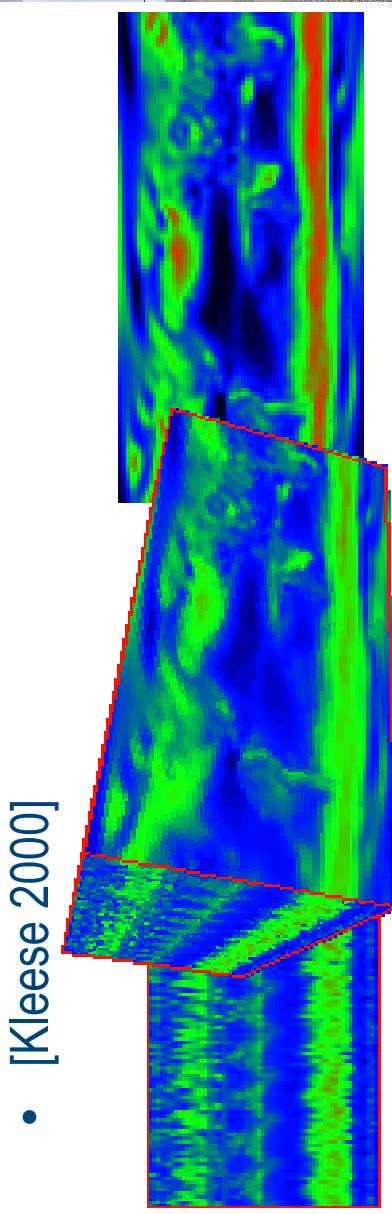
- 50+ physical parameters („variables“): temperature, wind speed x/y, humidity, pressure, CO<sub>2</sub>, ...

- 2.5 TB per variable

## ■ observation:

Huge volumes moved,  
**only part needed** (10:1)

- [Kleese 2000]



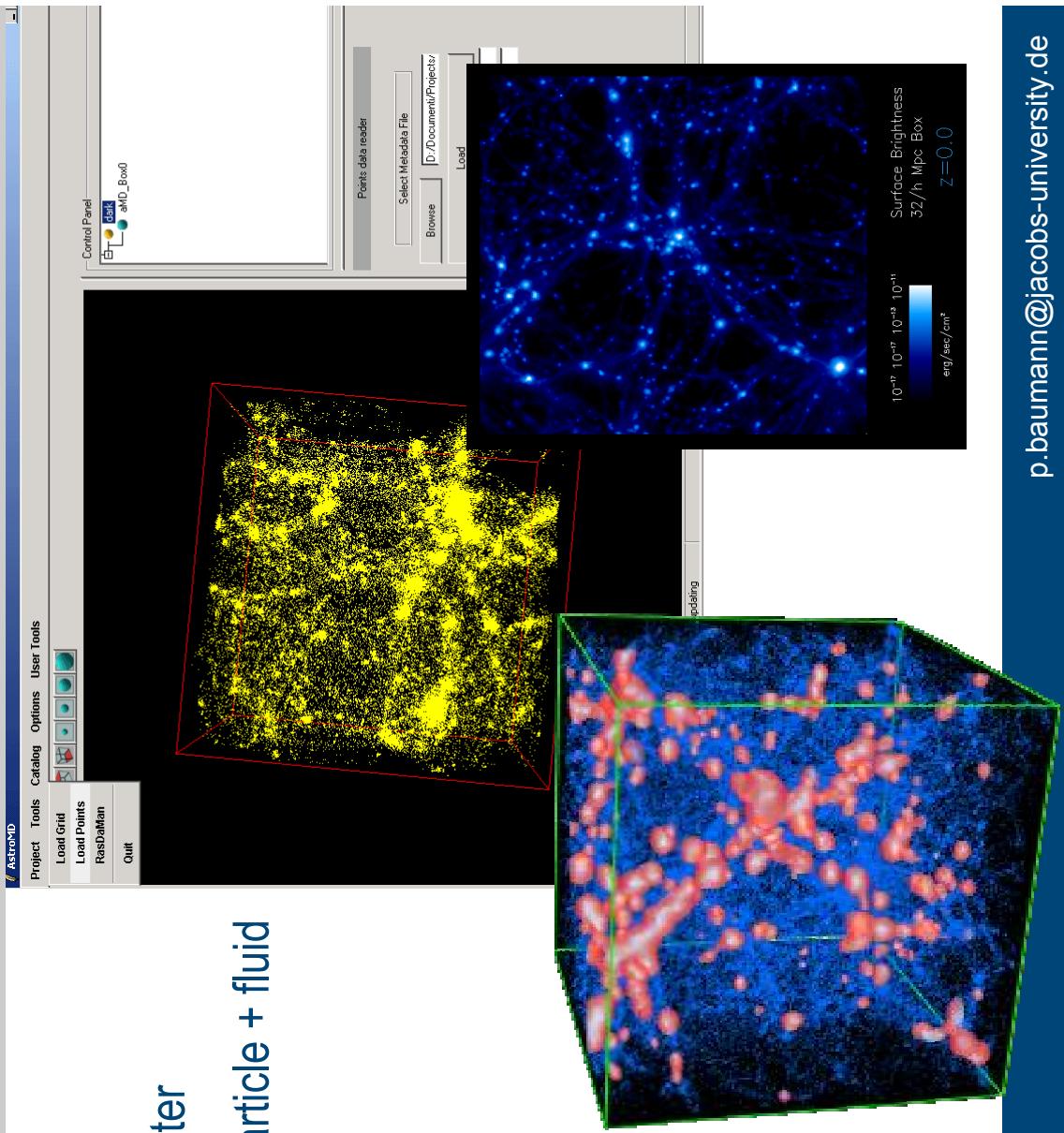
DKRZ: 24-node NEC SX-6



# Cosmological Simulation



- Modelling domain: 4-D
  - Dark matter, baryonic matter
  - → Coupled simulation: particle + fluid
- Results are 3-D/4-D cutouts from universe
  - Eg,  $64 \text{ Mpc}^3$   
(Mega Parsec;  
1 pc = 3.27 light years)



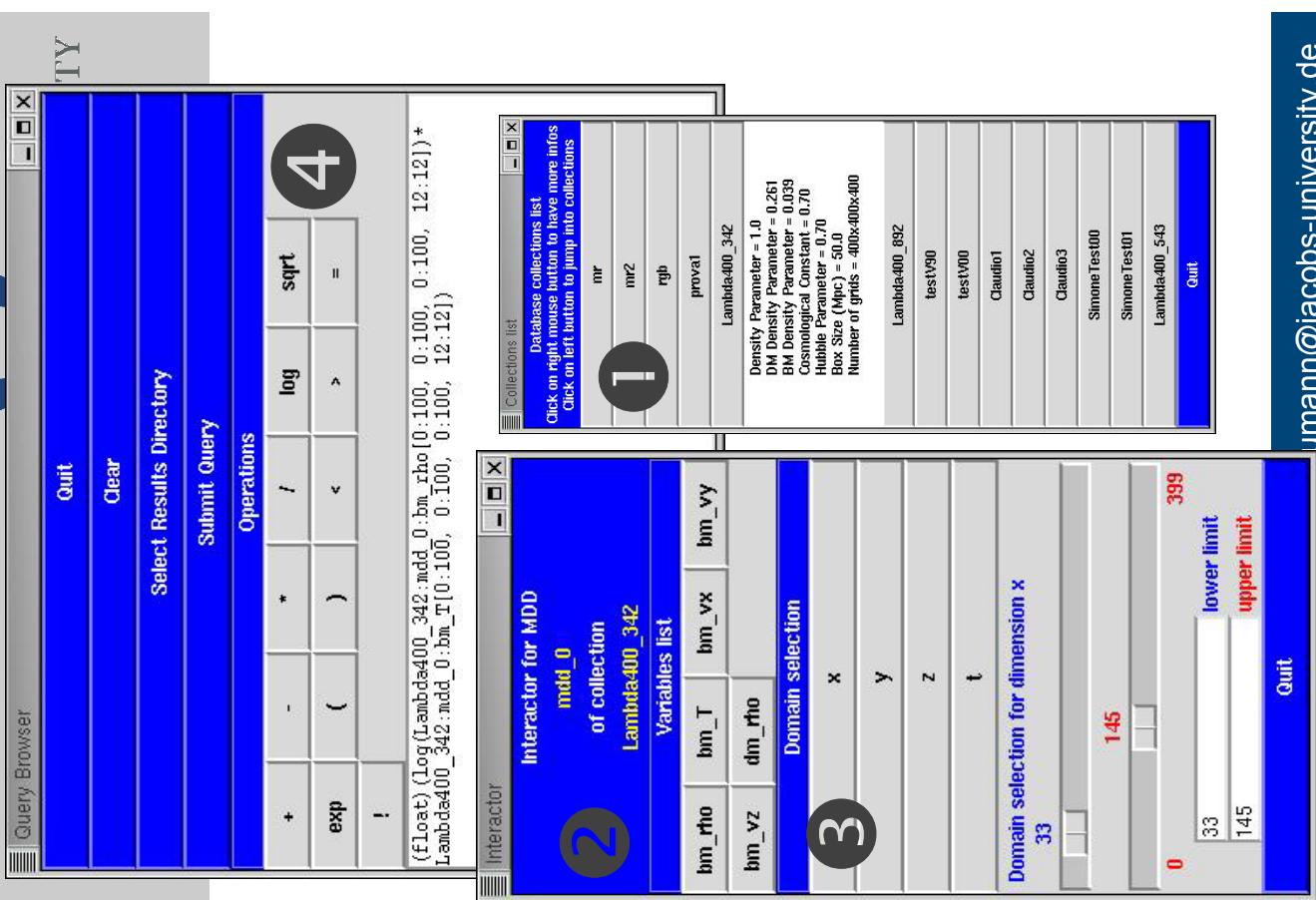
# Cosmology (contd.)

## ■ Guided retrieval:

- Selection of objects ①  
and their cell components ②
  - interactive setting of trim operations ③ per dimension ④
    - Augmented with induced operations ④

## ■ Suitable for expert users

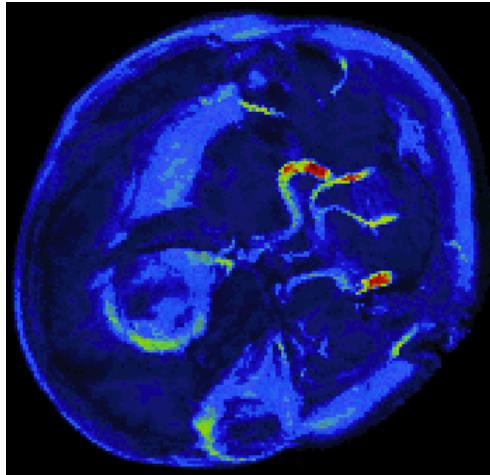
## ■ Details: [cosmolab.cineca.it](http://cosmolab.cineca.it)



# Human Brain Imaging



- Research goal: to understand structural-functional relations in human brain
- Experiments capture activity patterns (PET, fMRI)
  - Temperature, electrical, oxygen consumption, ...
  - → lots of computations → „activation maps“



- Example: “*a parasagittal view of all scans containing critical Hippocampus activations, TIFF-coded.*”

```
select tiff( ht [ $1, * : *, * : * ] )
from HeadTomograms as ht,
Hippocampus as mask
where count_cells( ht > $2 and mask )
/ count_cells( mask )
> $3
```

\$1 = slicing position, \$2 = intensity threshold value, \$3 = confidence

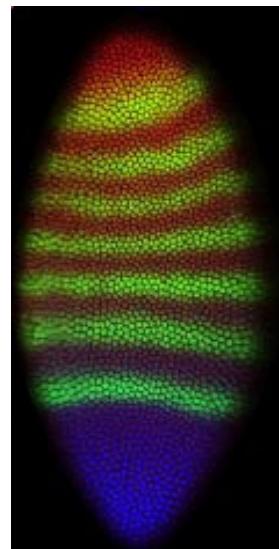
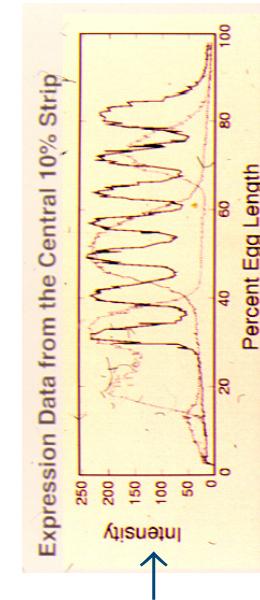
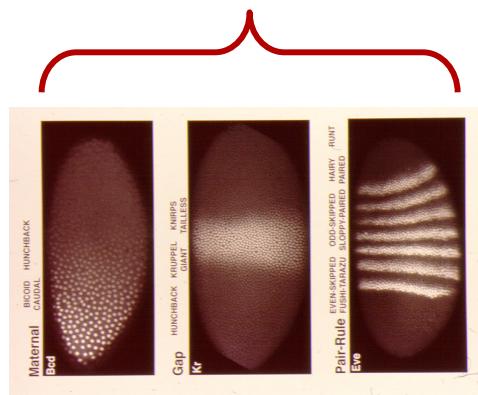
# Gene Expression Analysis



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[http://urchin.spbcas.ru/Mooshka/ \[Samsonova et al\]](http://urchin.spbcas.ru/Mooshka/)

- Gene expression = reading out genes for reproduction
- Research goal: capture spatio-temporal expression patterns in *Drosophila*



```
select jpeg( scale( { {1c, 0c, 0c}*e [0, * : *, * : *]  
+ {0c, 1c, 0c}*e [1, * : *, * : *]  
+ {0c, 0c, 1c}*e [2, * : *, * : *] , 0.2 ) )  
from EmbryoImages as e
```

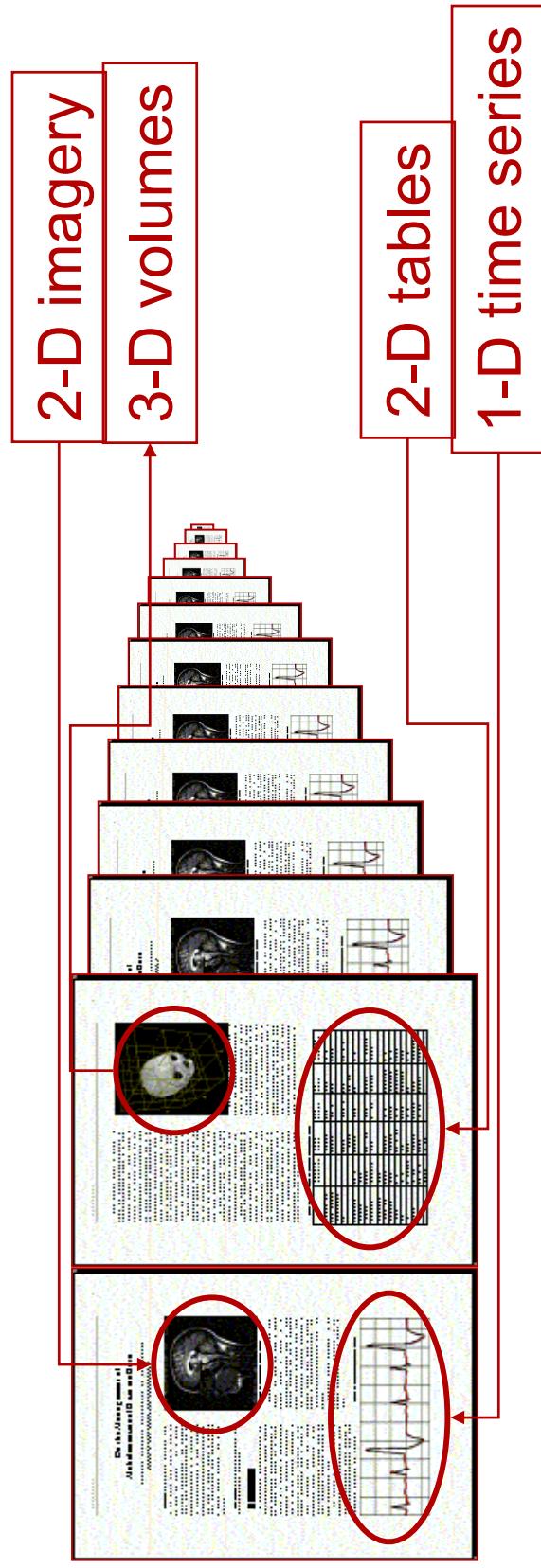
# Roadmap



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- Introduction
- Conceptual modelling
- Architecture
  - Arch I: Storage Management
  - Arch II: Query Processing
- Applications
- Wrap-up

# Vision: Document Integrated Retrieval



*„all clinical trials of drug X  
where patient temperature > 40° C within the first 48 hours.“*

# Finally...

- value-added raster data services important + growing field
  - Service providers & users demand it
  - Currently driven by geo apps
  - "2D, 3D imagery next great challenge in geo databases"  
[Xavier Lopez, Oracle]
- Many research issues in all facets
  - rasdaman System:  
commercialized + research vehicle
- contact:  
[p.baumann@jacobs-university.de](mailto:p.baumann@jacobs-university.de)

