# Example-Driven Design of Efficient Record Matching Queries 

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## Record Matching

## - Katrina: Given evacuee lists

match against enquiries.

|  | First Name | Last Name | Address | Phone | Father | Mother |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\rightarrow$ | Holmes | Elois | 2723 Third St | 938-8374 |  |  |
|  | Donneaka | Martin |  | 504-974-637 | Donald Qautier |  |
|  | Thomas |  | 2435 Delachise St |  | Lomax |  |
| $\square$ | Elois | Holmes | Third Street |  |  |  |
|  | Donaka | M |  | 504-974-637 | D Oautier |  |

## Record Matching can be Difficult

| First Name | Last Name | Address | Phone | Father | Mother |
| :--- | :--- | :--- | :--- | :--- | :--- |
| Holmes | Elois | 2723 Third St | $938-8374$ |  |  |
| Donneaka | Martin |  | $504-974-637$ | Donald Qautier |  |
| Thomas |  | 2435 Delachise St |  | Lomax |  |
| Elois | Holmes | Third Street |  |  |  |
| Donaka M  $504-974-637$ D Oautier  |  |  |  |  |  |

- Too many options to consider while building a record matching query
- Complicated due to errors and representational differences


## Record Matching Queries

```
Select * from Enquiries R, Evacuees S where
sim1(R.FirstName + R.LastName, S.FirstName + S.LastName) > 0.85
AND sim2(R.Address, S.Address) > 0.83
AND sim3(R.Phone, S.Phone) = 1
OR
sim4(R.FirstName + R.LastName + R.Phone,
    S.FirstName + S.LastName + S.Phone) > 0.87
OR
1.5 * sim5(R.FirstName, S.FirstName) - 0.3 * sim6(R.Father +
    R.Mother, S.Father + S.Mother) > 0.9
```


## Creating RM Queries

- Challenges
- Which column combinations to compare?

■ Which similarity function for each combination?

- Name similarity: soundex or edit distance
- Address similarity: jaccard
- How to determine the thresholds for chosen similarity function-column combination choice?


## Example-Driven Approach

- Input
- A set of example ( $\mathbf{r}, \mathrm{s}$ ) record pairs: matches \& non-matches
- A set of candidate operators
- Goal
- Construct a query which has the "best quality" when applied to the examples
- Quality measure
- Recall: Number of correctly classified matching pairs s.t. the fraction of false positives is less than $B$


## Previous Work

- Machine learning (ML) based predicates
- Decision trees
- SVMs - more accurate
- However, cannot efficiently implement similarity joins involving ML predicates
- Usually, cross product followed by filter


## SVM Predicates

- Current best method [Bilenko et al.]
- Example SVM predicate
- 1.5*Jaccard(R.[NAC], S.[NAC]) $0.3^{*} E d i t(R . N, S . N)>0.9$
- May not be efficiently executable
- Cross product followed by a filter


## Our Approach

- Constrain class of output queries
- Efficiently executable
- Flexible enough to capture a rich set of queries
- Programmers can review \& modify
- If required, add more sophisticated ML predicates to suggested queries


## Similarity Space

- Map examples to +/- points
- D-dimensional: One per similarity fn \& column combination
- Matches $\rightarrow$ + and Non-matches $\rightarrow$ -

Name similarity (edit similarity)
Address similarity (jaccard over ACZ)
Predicate: name similarity > c1 and address similarity >c2


## Class of Queries

- Relations R, S (schema [Name, Address, City, Zip])
- $D$ similarity functions (and column combinations)
- Class: Union of top-right rectangular boxes


## Similarity functions

Name similarity (edit similarity)
Address similarity (jaccard over ACZ)

name similarity > C1 and address similarity > c2
OR
name similarity > c1' and address similarity > C2'

## Problem Statement

- Given positive and negative points, find K rectangular boxes such that
- Recall-the number of positive points in them-is maximized
- Number of negative points they contain is less than $B$
- Sub-space constraints on each rectangular box
- Not more than d (<= D) dimensional


## Algorithm Outline

- Consider B=0
- No negative points at all in the result
- Extend to B>0
- Allow a few negative points in the result


## Union of Rectangles

- Find the best valid rectangular box with the maximum number of + 's
- Remove +'s in box and iterate



## Best Rectangular Box

- Recursive search for the best valid rectangular box

- Can be applied to $\mathrm{D}>2$ and for boxes in sub-spaces (i.e., d<D)


## Union of Rectangular Boxes

- Greedy strategy
- Pick best rectangular box with maximum number of +'s and no -'s
- Remove +'s contained in box
- Iterate until
- All +'s are covered
- K boxes are picked
- Approximation guarantee
- Within (1-1/e) of the optimal
- Follows from the greedy solution to the set coverage problem


## Allowing Non-matches

- A valid rectangular box may now include a fraction of negative points
- Find the best among all valid boxes
- Recursive algorithm applicable again



## Record Transformations

- Consider two records
- r1: [Matrin Smith, Redmond, WA, 98052]
- s1: [Martin Smit, NULL, WA, 98052]
- Apply FD zip $\rightarrow$ city to s1
- s1': [Martin Smit, Redmond, WA, 98052]
- For many similarity functions, $\operatorname{sim}(\mathrm{r} 1, \mathrm{~s} 1)$ < sim(r1, s1')
- Hence record transformations help identify matches!


## Record Transformations (contd)

- Example record transformations
- FDs to fill in missing values
- Splitting columns into sub-columns (e.g., address or product names)
- Our framework can be extended to consider such transformations
- Idea: Iteratively add best transformation to the current query


## Experimental Evaluation

- Datasets
- Organization data from an operational data warehouse
- RIDDLE repository ([Bilenko], UT Austin)

Techniques compared

- Addresses: a commercial cleansing tool called Trillium
- RIDDLE: SVM


## Operator Trees vs. Trillium

- 29 candidate similarity functions
- Zipcode splitter: out-code and in-code
- Out-code $\rightarrow$ City
- At most 4 similarity functions per box
- Union of at most 4 boxes

|  | Precision | Recall |
| :---: | :---: | :---: |
| Trillium | 0.99 | 144 K |
| Operator Trees | 0.98 | 159 K |
| Baseline | 0.98 | 80 K |

## Cora Dataset

## - Bibliography data: authors, titles



## Efficiency of Similarity Join

Similarity join (jaccard similarity) over 500K record relation with itself

- [VLDB06] SSJoin algorithm

| Threshold | SimJoin |
| :---: | :---: |
| 0.9 | 61 s |
| 0.85 | 125 s |
| 0.80 | 285 s |

SVM predicate: 10 days
SVM + blocking: 1+ hour

## Conclusions

- Example-driven approach to suggest a record matching query
- Considered constrained space of efficiently executable queries
- Empirically demonstrated accuracy
- Web search: "data cleaning project"
- http://research.microsoft.com/dmx/datacleaning


## Questions

