Stop and Restart Style Query Execution

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Motivation

- Long running decision support queries can be resource intensive
 - Resource contention can be reduced by terminating queries
- Terminated queries have to be re-rerun completely
 - Irrespective of how close to completion they were
- Can we do better?
 - Assume database is not updated





Problems with this Approach

- Overheads can be unbounded
 - e.g. Filter predicate selects most of the records
 - Need to cache results in memory or periodically flush them to disk
 - The query may not be terminated!
- Bounded overhead
 - Save and Reuse "Best" K records
 - Flush results to disk when the query is terminated

Outline

Primitives

- Bounded Query Checkpointing
- Opt-Skip Algorithm
- Complex Query Plans
- Experiments

Skip-Scan Operator

- Generalization of Scan Operator
- Takes 2 parameters (LB, UB)
- Skips all records between LB and UB
- Sends a EOLB message after processing LB





Cost Model



Work done by query = Total number of getNext Calls issued across all operators



Bounded Query Checkpointing

Parameters

- Budget k
- Let id be the current ID of the record in the leaf node
- Among all bounded RPlan(LB,UB) where LB < UB <= id</p>
 - Maintain the restart plan with maximum benefit

Candidate Windows and Restart Plans (k=3)



- Track candidate windows (r_{i-1},r_{i+k}) of size k+2
 - Derive LB and UB values from the source RIDS corresponding to r_{i-1} and r_{i+k}
- Let GN(ri) denote total GetNext calls issued in the pipeline until r_i was generated at the pipeline root

Benefit of Restart Plan = GN(r_{i+k}) – GN(r_{i-1})

Opt-Skip Algorithm

/* window, k= total budget */
/* BestW = best window */

```
Algorithm Opt-Skip
BestW = empty set
W = empty set
For Each intermediate record r_i do:
Append r_i to W
If W.Size() > k+2 then
W = last k+2 records in W
SkippableW = FindSkippable(W)
```

```
If Benefit(SkippableW) > Benefit(BestW) then
BestW = SkippableW
```

Restart plans for Multiple Pipelines



Algorithms for Multiple Pipelines

- Current Pipeline
 - Use Opt-Skip on the current pipeline
- Max-Pipeline
 - Choose single "best" pipeline from all completed pipelines
- Merge-Pipeline
 - Heuristic to distribute the budget k among multiple pipelines
- Subtree Caching

Experimental Evaluation

- Prototyped in Microsoft SQL Server 2005
 - Built skip-scan operator on top of the clustered index scan operator

Evaluation Metric

- T1 = Total number of GetNext calls in Initial Run before termination
- T2 = GetNext calls issued in the restart to reach the same point in execution
- Percentage Work Saved (PWS) = (T1-T2)/T1*100

Goals

Effect of clustering

- Utility of bounded query checkpointing for complex queries
- Overheads in the initial run
- Algorithms for multi pipeline plans



Effect of Clustering



TPC-H Queries (Termination at 50%)



Overheads in Initial Run were < 3% for most Queries

Multi-Pipeline Algorithms



SkyServer Database



Summary of Results

- Many cases where bounded checkpointing is useful
 - selective predicates, correlation with the clustering column, subtree caching
- Saving a small number of records can result in substantial savings
- Keeping track of previous pipelines is important

Future Work

- Aggregations
 - Saving partial sums
- Hash Spills
 - Maintaining the skipping invariant
- Handling Updates
 - Validate Restart Plan

Related Work

SQL Cursors

Does not release any state

Query Resumption

- [Labio et al. SIGMOD 2000]
- [Chandramouli et al. SIGMOD 2007]
- Delta results vs. Full Results
- Skip-Scan can be utilized for resuming the build phase of hybrid hash join

Conclusions

Resource contention can lead to termination of queries

- Queries have to be rerun completely
- Bounded Query Checkpointing
 - Skip-Scan Operator
- Saving a small amount of records can lead to substantial savings